

Wang (2002) introduced a wide range of heuristic search algorithms. However, the performance of these algorithms is highly dependent on the problem. The performance of the genetic algorithm (GA) is highly dependent on the problem. The performance of the genetic algorithm (GA) is highly dependent on the problem. The performance of the genetic algorithm (GA) is highly dependent on the problem.

Other heuristic search algorithms (PTL) (Draetta 1994; Draetta 1996, 2006) are also used for optimization. The performance of these algorithms is highly dependent on the problem. The performance of these algorithms is highly dependent on the problem. The performance of these algorithms is highly dependent on the problem.

The performance of the heuristic search algorithms is highly dependent on the problem. The performance of these algorithms is highly dependent on the problem. The performance of these algorithms is highly dependent on the problem. The performance of these algorithms is highly dependent on the problem.

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2 p - π

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$$\pi ::= x(y) \mid \bar{x}(y) \mid \tau \mid I_p \mid skip \mid \varepsilon$$

$$P ::= 0 \mid \pi \cdot P \mid P_1 + P_2 \mid P_1 \mid P_2 \mid [a_1 = a_2]P \mid \text{va } P \mid A\langle a_1, \dots, a_n \rangle$$

where x, y, a, a_i (i is a integer $1 \leq n \in \mathbb{N}$) are variables and \mathcal{N} . The action π is a sequence of actions, and \mathbb{N} is the set of natural numbers. The action π is a sequence of actions, and \mathbb{N} is the set of natural numbers.

ceci ig (de ed b x(y)) e age, e ec i ga ie τ sa ii ε .
 A i α a aci e e e ei ha $I_p = \{p_1 \wedge skip, \dots, p_k \wedge skip\}$ e $skip$.
 Ha e $p_i \in PR (1 \leq i \leq k \text{ a } d_i \in \mathbb{N})$, a d $skip$ i a ecia s i i dica i ga
 i e, i. F s c e ie ce w e, e \tilde{p}_i e e e $p_i \wedge skip$ a d $\bar{x}(x)$ de e
 $\bar{x}()$ (s x()). N e ha $\bar{x}()$ (s x()) e a a aci e w i h, e age.

A s ce (s s ce e e i) ca be a s ce 0, a aci e
 g aded s ce $\pi \cdot P$, ai f_w s ce e $P_1 + P_2$, a e e i i
 f_w s ce e $P_1 | P_2$, a ch s, c, e $[a_1 = a_2]P$, a e s i c i s, c, e va P , s
 a s ce i a ce $A(a_1, \dots, a_n)$. $[a_1 = a_2]P$ i dica e ha P i e ec ed if $a_1 = a_2$.
 va P beha e a P e ce ha he s, i ca i he b, d e a i f s-
 bidde. A i $A(a_1, \dots, a_n)$ i a s ce ide i α de ed b $A(x_1, \dots, x_n) \stackrel{\text{def}}{=} P_A$
 a d $A(a_1, \dots, a_n) = \{\vec{a} / \vec{x}\} P_{A_w}$ ha e \vec{a} a d \vec{x} a e he ec s f_{a_1}, \dots, a_n a d
 x_1, \dots, x_n , e ecie. The e f e i a s ce $n(P)$ c i fb, d e
 $bn(P)$ a d f e e $fn(P)$. N e, a a, a e a i g i (a) s va P a e b, d
 e. O ha a e f e e. The ab e i a i ab, $skip$ a d $await$ a e de ed
 a f w .

$$\begin{aligned} \text{Skip}_0 & \quad skip^0 \stackrel{\text{def}}{=} \varepsilon \\ \text{Skip}_n & \quad skip^n \stackrel{\text{def}}{=} skip \cdot skip^{n-1} (n \geq 1) \\ \text{Await}_d \text{ await}(P) & \stackrel{\text{def}}{=} \varepsilon \cdot P + skip \cdot P + \dots + skip^n \cdot P (n \geq 0) \end{aligned}$$

w ha e $n \in \mathbb{N}$. The da i ed s ce $await(P)$ i, ed s e a i e ch, s-
 i, i ca i. Ha e P i $await(P)$ i a e di g (s e ceci ig) e g aded s ce.

Definition 1 The s, c, s a c g, e ce α he e f π s ce e \mathcal{P} i de ed b
 he i e e, a i be w .

- S1 $P|0 \equiv P$ S2 $\nu x P \equiv \nu y \{y/x\}P$ if $y \notin fn(P)$
- S3 $P + Q \equiv Q + P$ S4 $x(y) \cdot P \equiv x(z) \cdot \{z/y\}P$ if $z \notin fn((y)P)$
- S5 $P|Q \equiv Q|P$ S6 $(P|Q)|R \equiv P|(Q|R)$
- S7 $\nu x 0 \equiv 0$ S8 $\nu x (P|Q) \equiv P|\nu x Q$, if $x \notin fn(P)$
- S9 $\nu xy P \equiv \nu yx P$ S10 $[x = y]P \equiv 0$, if $x \neq y$
- S11 $\varepsilon \cdot P \equiv P$ S12 $skip \cdot P + skip \cdot Q \equiv skip \cdot (P + Q)$
- S13 $P + 0 \equiv P$ S14 $A(a_1, \dots, a_n) \equiv \{\vec{a} / \vec{x}\}P_A$ if $A(x_1, \dots, x_n) \stackrel{\text{def}}{=} P_A$
- S15 $\nu x \bar{x}(y) \cdot P \equiv 0$ S16 $P + (Q + R) \equiv (P + Q) + R$

The ec i f π s ce e c i f_w age: e f s i a a e, aci
 e e a d he ha e f s i α a aci e e a d he e c i f a i α a
 aci e i a e d a e a i α a w i h e, i. F s he α a i a
 e, he b α a e a c i i g i e a f w :

$$\begin{aligned} \pi_o & ::= \bar{x}(y) | x(z) | \bar{x}(z) | \tau | I_p | skip \\ \pi_c & ::= \bar{x}(y) | x(z) | \bar{x}(z) | \tau \\ \pi_t & ::= I_p | skip \end{aligned}$$

<p>Tau: $\frac{}{\tau \cdot P \xrightarrow{\tau} P}$</p> <p>Out: $\frac{}{\bar{x}(y) \cdot P \xrightarrow{\bar{x}(y)} P}$</p> <p>Sum: $\frac{P_1 \xrightarrow{\pi_c} P'_1}{P_1 + P_2 \xrightarrow{\pi_c} P'_1}$</p> <p>Com: $\frac{P_1 \xrightarrow{\bar{x}(y)} P'_1 \quad P_2 \xrightarrow{x(z)} P'_2}{P_1 P_2 \xrightarrow{\tau} P'_1 \{y/z\} P'_2}$</p> <p>Open: $\frac{P \xrightarrow{\bar{x}(y)} P'}{\nu y P \xrightarrow{\bar{x}(z)} \{z/y\} P'} \quad (x \neq y, z \notin fn(\nu y P'))$</p> <p>Res: $\frac{P \xrightarrow{\pi_o} P'}{\nu x P \xrightarrow{\pi_o} \nu x P'} \quad (x \notin n(\pi_o))$</p> <p>Act_t: $\frac{}{\pi_t \cdot P \xrightarrow{\pi_t} P}$</p> <p>Act_ε: $\frac{P \xrightarrow{\pi_c} P'}{\varepsilon \cdot P \xrightarrow{\pi_c} P'}$</p> <p>Com_{idle}: $\frac{P_1 \xrightarrow{\pi_t} P'_1 \quad P_2 \equiv 0}{P_1 P_2 \xrightarrow{\pi_t} P'_1}$</p> <p>Com_t: $\frac{P_1 \xrightarrow{\pi_{t_1}} P'_1 \quad P_2 \xrightarrow{\pi_{t_2}} P'_2}{P_1 P_2 \xrightarrow{\pi_{t_1} \cup \pi_{t_2}} P'_1 P'_2} \quad (\pi_{t_1} \cap \pi_{t_2} - skip = \emptyset, P_1 P_2 \not\xrightarrow{\tau})$</p>	<p>In: $\frac{}{x(z) \cdot P \xrightarrow{x(w)} \{w/z\} P} \quad (w \notin fn((z)P))$</p> <p>Par: $\frac{P_1 \xrightarrow{\pi_c} P'_1}{P_1 P_2 \xrightarrow{\pi_c} P'_1 P_2} \quad (bn(\pi_c) \cap fn(P_2) = \emptyset)$</p> <p>Close: $\frac{P_1 \xrightarrow{\bar{x}(y)} P'_1 \quad P_2 \xrightarrow{x(y)} P'_2}{P_1 P_2 \xrightarrow{\tau} \nu y (P'_1 P'_2)}$</p> <p>Mat: $\frac{P \xrightarrow{\pi_o} P'}{[a_1 = a_2] P \xrightarrow{\pi_o} P'} \quad (a_1 = a_2)$</p> <p>Ide: $\frac{\{\bar{a}'/\bar{x}\} P_A \xrightarrow{\pi_o} P'}{A(\bar{a}') \xrightarrow{\pi_o} P'} \quad (A(\bar{x}) \stackrel{\text{def}}{=} P_A)$</p> <p>Sum_{t1}: $\frac{P_1 \xrightarrow{skip} P'_1 \quad P_2 \xrightarrow{skip} P'_2}{P_1 + P_2 \xrightarrow{skip} P'_1 + P'_2}$</p> <p>Await: $\frac{\varepsilon \cdot P + skip \cdot await(P) \xrightarrow{\pi_o} P'}{await(P) \xrightarrow{\pi_o} P'}$</p> <p>Sum_{t2}: $\frac{P_1 \xrightarrow{\pi_t} P'_1}{P_1 + P_2 \xrightarrow{\pi_t} P'_1} \quad (P_1 + P_2 \not\xrightarrow{skip})$</p>
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Fig. 1 Operational semantics of π

\mathcal{W} has a dyadic free variable $i \in fn(\pi_o)$ and $z \in fn(\pi_c)$. z is a bound variable, defined by $bn(\pi_o)$ and $bn(\pi_c)$. Here, $\bar{x}(z)$ is defined by $\nu z \bar{x}(z)$.

As in [12], a state i is a configuration $P \xrightarrow{\pi_o} Q$ if P can evolve to Q via π_o . In Fig. 1, the operational semantics of π is defined.

Definition 2 For a π -LTS P , the LTS (Labelled Transition System) of P is a pair (P, \mathcal{T}) , where \mathcal{T} is the set of transitions of P . For a π -LTS P , the LTS of P is defined as follows:

1. $P \in \mathcal{P}$;
2. For a $Q \in \mathcal{P}$, if $Q \xrightarrow{\pi_o} Q'$, then $Q' \in \mathcal{P}$, $(Q, \pi_o, Q') \in \mathcal{T}$. For a π -LTS P , the LTS of P is defined as follows:

3 Framed temporal logic programming language MSVL

Our design goal is to provide a logic (PTL) and MSVL in a declarative way, based on PTL. MSVL is a declarative logic. The aim of this section is to define the semantics of MSVL as follows:

$$e ::= c|v|v_p|\&x|*v_p|\bigcirc x|\bigcirc x|e_0ope_1 \quad (op ::= +|-|*|/)$$

$$b ::= \text{true}|\text{false}|\neg b|b_0 \wedge b_1|e_0 = e_1|e_0 < e_1$$

w ha e c i a c a , v a d a a a i a b e , v p a i a a i a b e a d x d e e a d a a a i a b e s a i a a i a b e . & i a i a e f a e c e a d * a i a d a e f a e c e . \bigcirc e a n e x t a d \bigcirc p r e v i o u s . A d p i c a i a b e x i a i d b e f u e d i s g p p r o g i f f r a m e (x) s l b f (x) i c a i e d i p r o g .

A f u e d s g p i M S V L c a b e f u a i e d b h e i e e e p e e a a p e b e $w . p_1, \dots, p_m, p$ a d q a e g e a a f u e d s g p .

T a i a i :	ε	B a i c a i g u e :	$x = e$
P i a a i g u e :	$*v_p = e$	S a e f u e :	$\text{lbf}(x)$
I a a f u e :	$\text{frame}(x)$	C j c i :	$p \wedge q$
S e e c i :	$p \vee q$	N e a p e :	$\bigcirc p$
A w a a p e :	$\square p$	C d i i a a p e :	$\text{if } b \text{ then } p$
e l s e q			
E i e i a , a i c a i :	$\exists x : p(x)$	S e , e i a a p e :	$p ; q$
P a a e :	$p \parallel q$	W h i e a p e :	$\text{while } b \text{ do } p$
S c h i e d c u i , i c a i :	$\text{await}(b)$	R j e c i :	$(p_1, \dots, p_m) \text{Prj } q$

ε i h e a i a i a p e . T h e b a i c a i g u e $x = e$ e a h a h e a , e f a i a b e x i e , a h e a , e f e s e i e . S i i a , * v p = e i h e a i g u e a c i a e d w i h h e i a . T h e , i a i g u e $x := e$ i d e d e d a $x := e \stackrel{\text{def}}{=} \text{skip} \wedge \bigcirc x = e$. T h e n e x t a p e $\bigcirc p$ e a h a p h d a h e i i e d i a e , c c e s a e . $\square p$ i e h a p h d i a a e f u w . T h e e , e i a a p e p ; q i g i e h a p h d f i h e c c e a e , i i e i i f i s e a w h i c h i a e a d q w i a e e c i g f u h a i .

T h e a e f u e l b f (x) a d i a a f u e f r a m e (x) , i i e f u i g e c h i , e . T h e f u i g a s i d e e d b f r a m e a d f r a m e (x) e a h a v a r i a b l e x a l w a y s k e e p s i t s o l d v a l u e o v e r a n i n t e r v a l i f n o a s s i g n m e n t t o x i s e n c o u n t e r e d . T g i e i d e i i , a w a i g u e i f u a i e d a $x \leftarrow e \stackrel{\text{def}}{=} x = e \wedge p_w$ h a e p x i a a i c s i i c e c e d w i h a i a b e x a d c a b e , e d f s h a s e . T h e h e a i g u e a g i f u a i e d a : a f (x) \stackrel{\text{def}}{=} p_x . T h , l b f (x) a d f r a m e (x) c a b e d e d a l b f (x) \stackrel{\text{def}}{=} \neg a f (x) \rightarrow \exists a : (\bigcirc x = a \wedge x = a) a d f r a m e (x) \stackrel{\text{def}}{=} \square (\text{more} \rightarrow \bigcirc \text{lbf}(x)) s e e c i e w h a e a i a a i c a i a b e .

T h e c d i i a a p e i f b t h e n p e l s e q s e a , a e b ; i f b i t r u e , h e p i e e c e d , h a w i e q i e e c e d . T h e i a i w h i l e b d o p a w s c e p b e e e a d e e c e d a i e (s i i e) b a f u e a g a h e c d i i b i a i e d a h e b e g i i g f e a c h e e c i . T h e e e c i a p e p \vee q s e e e h a p s q w i b e e e c e d . T h e c j c i a p e p \wedge q d e c a e h a h e s c e e p a d q a e e e c e d c c e e h a i g a h e a e a d a i a b e d i g h a , a e e c i . T h e a e c s c i p \parallel q , h w a h a c c e c u , a i i a a . T h e d i i g i h e d d i f f a e c e b e w e e p \parallel q a d p \wedge q i h a h e f u a w b h p a d q b e a b e e c i f h e i w i a w h i e h e a a d e . E . g . , l e n (2) \parallel l e n (3) h d b l e n (2) \wedge l e n (3) i b i ,

For each formula φ we consider its \mathcal{F} -reduct $\varphi|_{\mathcal{F}}$. According to the definition of $n(P)$ and its \mathcal{F} -reduct $n(P)|_{\mathcal{F}}$, we have $\text{CH}_{\mathcal{F}(P)}$ and $\text{MSVL}_{\mathcal{F}(P)}$ are based on \mathcal{F} . For instance, if $n(P) = \{a, b, c\}$ and $\text{PRP} = \{p_1, p_2, p_3\}$, we have $\text{CH}_{\mathcal{F}(P)} = \{C_a, C_b, C_c\}$ and $V_{\mathcal{F}(P)} = \{x_{p_1}, x_{p_2}, x_{p_3}\}$. The \mathcal{F} -reduct of a formula φ is defined by $\varphi|_{\mathcal{F}}$.

4.1 \mathcal{F} -reduct

$$\mathcal{F}(0) \stackrel{\text{def}}{=} \varepsilon$$

\mathcal{F} -reduct of a digraph, is a graph G such that $G|_{\mathcal{F}}$ is MSVL.

4.2 Order-acyclic

$$\mathcal{F}(\bar{x}(y) \cdot P) \stackrel{\text{def}}{=} \text{e } d(C_x, C_y); \mathcal{F}(P)$$

Order-acyclic graphed sets $\bar{x}(y) \cdot P$ are digraphs, where \bar{x} and y behave as P . Since $\bar{x}(y)$ is a \mathcal{F} -reduct, order-acyclic graphed sets $\bar{x}(y) \cdot P$ are \mathcal{F} -reducts of $d(C_x, C_y); \mathcal{F}(P)$. Hence, $\mathcal{F}(\bar{x}(y) \cdot P)$ is a \mathcal{F} -reduct of $d(C_x, C_y)$ and \mathcal{F} is MSVL-cha $\bar{x}(y) \cdot P$.

4.3 In-acyclic

$$\mathcal{F}(x(y) \cdot P) \stackrel{\text{def}}{=} \exists C_y : (\text{frame}(C_y) \wedge \text{e } e(C_x, C_y); \mathcal{F}(P))$$

In-acyclic graphed sets $x(y) \cdot P$ indicate \mathcal{F} -reducts $\bar{x}(y) \cdot P$ and the behavior $\{z/y\}P$. Hence, $\bar{x}(y) \cdot P$ is a \mathcal{F} -reduct of $\bar{x}(y) \cdot P$. The \mathcal{F} -reduct of $\bar{x}(y) \cdot P$ is $\bar{x}(y) \cdot P|_{\mathcal{F}}$. For $\bar{x}(y) \cdot P$, $\text{frame}(C_y)$ is a \mathcal{F} -reduct of $\bar{x}(y) \cdot P$ and $\bar{x}(y) \cdot P|_{\mathcal{F}}$ is a \mathcal{F} -reduct of $\bar{x}(y) \cdot P|_{\mathcal{F}}$. What we have here is that C_x is a \mathcal{F} -reduct of $\bar{x}(y) \cdot P$ and C_y is a \mathcal{F} -reduct of $\bar{x}(y) \cdot P$. We have $\bar{x}(y) \cdot P|_{\mathcal{F}}$ is a \mathcal{F} -reduct of $\bar{x}(y) \cdot P|_{\mathcal{F}}$. As a graph, $\bar{x}(y) \cdot P$ is a \mathcal{F} -reduct of $\bar{x}(y) \cdot P|_{\mathcal{F}}$.

4.4 \mathcal{F} -await

$$\mathcal{F}(\text{await}(\bar{x}(y) \cdot P)) \stackrel{\text{def}}{=} \text{ai } _ e d(C_x, C_y); \mathcal{F}(P)$$

$$\mathcal{F}(\text{await}(x(y) \cdot P)) \stackrel{\text{def}}{=} \exists C_y : (\text{frame}(C_y) \wedge \text{ai } _ e(C_x, C_y); \mathcal{F}(P))$$

For $\text{await}(P)$, we have P is a digraph (acyclic) graphed set, where $\bar{x}(y) \cdot P$ is a \mathcal{F} -reduct of $\bar{x}(y) \cdot P$. The \mathcal{F} -reduct of $\bar{x}(y) \cdot P$ is $\bar{x}(y) \cdot P|_{\mathcal{F}}$. Hence, $\bar{x}(y) \cdot P|_{\mathcal{F}}$ is a \mathcal{F} -reduct of $\bar{x}(y) \cdot P|_{\mathcal{F}}$. As a graph, $\bar{x}(y) \cdot P$ is a \mathcal{F} -reduct of $\bar{x}(y) \cdot P|_{\mathcal{F}}$.

sa fwait $(\bar{x}\langle y \rangle \cdot P)$ a dawait $(x\langle y \rangle \cdot P)$ a eci e .

4.5 I a aci e

$$\mathcal{F}(\tau \cdot P) \stackrel{\text{def}}{=} p_I \wedge \varepsilon; \mathcal{F}(P)$$

I a aci e gaded s ce $\tau \cdot P$ a fwait a aci e a d he beha e a P . I ca b a ed $p_I \wedge \varepsilon; \mathcal{F}(P)$ i w hich p_I i a eci e s i i ed a i f e w he h a s a i a aci a e ace.

4.6 R a aci e

$$\mathcal{F}(I_p \cdot P) \stackrel{\text{def}}{=} x_{p_1} = \text{true} \wedge \dots \wedge x_{p_n} = \text{true} \wedge \text{skip}; \mathcal{F}(P)$$

w h a e $I_p = \{\tilde{p}_1, \dots, \tilde{p}_n\}$. The a aci e gaded s ce $I_p \cdot P$ a i e $p_i (1 \leq i \leq n)$ i he s i e, i a d he beha e a P . I i e, i ca b a ed $x_{p_1} = \text{true} \wedge \dots \wedge x_{p_n} = \text{true} \wedge \text{skip}; \mathcal{F}(P)$. H a e, $x_{p_1}, \dots, a d x_{p_n}$ a e h e c s e d i g b e a a i a b e f $p_1, \dots, a d p_n$, s e eci e .

4.7 i e, i aci e

$$\mathcal{F}(\text{skip} \cdot P) \stackrel{\text{def}}{=} \text{skip}; \mathcal{F}(P)$$

i e, i aci e gaded s ce $\text{skip} \cdot P$ d e a e h a i w i d e i h e s i e, i a d he beha e a P . I s a fwait i s a i g h f w a d.

4.8 E aci e

$$\mathcal{F}(\varepsilon \cdot P) \stackrel{\text{def}}{=} \varepsilon; \mathcal{F}(P)$$

F a e aci e gaded s ce $\varepsilon \cdot P$, i w i e e c e a f s a i i a d he beha e a P . Si i a $\text{skip} \cdot P$, i i d i e c s a a e d i $\varepsilon; \mathcal{F}(P)$.

4.9 N d e i i c h i c e

$$\mathcal{F}(P_1 + P_2) \stackrel{\text{def}}{=} \mathcal{F}(P_1) \vee \mathcal{F}(P_2)$$

S i c s e $P_1 + P_2$ h w h a P_1 a d P_2 i s c e e d d e a i a e , h a i c a a s a b e e s e e d b a d i j, c i a e .

4.10 Parallels

$$\mathcal{F}(P_1 \parallel P_2) \stackrel{\text{def}}{=} \mathcal{F}(P_1) \parallel \mathcal{F}(P_2)$$

Parallels process $P_1 \parallel P_2$ is a process where P_1 and P_2 execute in parallel. A definition of parallel composition is given below.

4.11 Match

$$\mathcal{F}([a_1 = a_2]P) \stackrel{\text{def}}{=} \text{if } C_{a_1} = C_{a_2} \text{ then } \mathcal{F}(P) \text{ else } \varepsilon$$

Match process $[a_1 = a_2]P$ is a process where P is executed only if a_1 and a_2 are equal. If a_1 and a_2 are not equal, the process terminates. A definition of match is given below.

4.12 Recursive

$$\mathcal{F}(va P) \stackrel{\text{def}}{=} \exists C_a : (\text{frame}(C_a) \wedge \mathcal{F}(P))$$

Recursive process $va P$ is a process where P is executed with a fresh name a . The process $va P$ is defined as $\exists C_a : (\text{frame}(C_a) \wedge \mathcal{F}(P))$. A definition of recursive process is given below.

4.13 Recursive

Recursive process $A\langle a_1, \dots, a_n \rangle_{\mathcal{W}}$ is a process where $A(x_1, \dots, x_n)$ is defined by P_A , where $A\langle a_1, \dots, a_n \rangle = \{\vec{a} / \vec{x}\} P_A$. The process $A\langle a_1, \dots, a_n \rangle_{\mathcal{W}}$ is defined as $\mathcal{F}(va P)$.

Theorem 1 Let process $A\langle x_1, \dots, x_n \rangle \stackrel{\text{def}}{=} P_A$ and $A\langle a_1, \dots, a_n \rangle = \{\vec{a} / \vec{x}\} P_A$. If A is recursively called in P_A and there is a P'_A in the form of $\pi_1^s + \dots + \pi_i^s + \pi_1'^s \cdot A\langle x_1, \dots, x_n \rangle + \dots + \pi_m'^s \cdot A\langle x_1, \dots, x_n \rangle$ such that $A\langle x_1, \dots, x_n \rangle \equiv P'_A$, where, π_j^s ($1 \leq j \leq i$) and $\pi_k'^s$ ($1 \leq k \leq m$) denote $\pi_{j_1} \dots \pi_{j_n}$, $\pi_{k_1} \dots \pi_{k_n}$, respectively, then, there exists the least fixed point $L_{fix}\langle a_1, \dots, a_n \rangle$ of $A\langle a_1, \dots, a_n \rangle$ such that $A\langle a_1, \dots, a_n \rangle = L_{fix}\langle a_1, \dots, a_n \rangle$.

Proof The theorem can be proved by using the induction method (Tay 1955) and Scott's induction principle (Wierwille 1993). \square

Finally, we define the recursive process $A\langle x_1, \dots, x_n \rangle = \pi_1^s + \pi_1'^s \cdot A\langle x_1, \dots, x_n \rangle$, where π_1^s and $\pi_1'^s$ are defined as $\pi_1^s = \{\vec{a} / \vec{x}\} \pi_1'^s \cdot \pi_1^s$. Here, π_1^s and $\pi_1'^s$ are defined as $\pi_1^s = \{\vec{a} / \vec{x}\} \pi_1'^s \cdot \pi_1^s$. The recursive process $A\langle a_1, \dots, a_n \rangle$ is defined as $L_{fix}\langle a_1, \dots, a_n \rangle$.

MSVL_w e ca ecif a abis a i eg a N a b a fi a ai i e c s
 hes ec s i . F s A(x₁, ..., x_n) = π₁^s + ... + π_i^s + π₁^s · A(x₁, ..., x_n) + ... + π_m^s ·
 A(x₁, ..., x_n), he s a f ai i i a . Si ce a s ce i a ce ca be de ed
 w i h e c s i e ca i -π, i s a f ai w i be di ided i h ee ca e :

1. Wi h , s ec s i e ca i he de i i f A:

$$\mathcal{F}(A(a_1, \dots, a_n)) \stackrel{\text{def}}{=} \{\vec{C}_a / \vec{C}_x\} \mathcal{F}(P_A)$$

F s hi i d f s ce i a ce, i w i be di ec s a a ed b , b i , i . \vec{C}_a
 a d \vec{C}_x a e he c s e di g MSVL cha e ec s f \vec{a} a d \vec{x} . H a e w e , e
 a , b i , i g e h d $\{\vec{C}_a / \vec{C}_x\}$, b i , e \vec{C}_a f s \vec{C}_x ha each C_{xw} i be
 , b i , ed b C_{aw} i hi he b , d c e f C_{x_i} (1 ≤ i ≤ n) i $\mathcal{F}(P_A)$.

2. Wi h e c s i e ca i A a d A(x₁, ..., x_n) ≡ π₁^s + π₁^s · A(x₁, ..., x_n):

$$\begin{aligned} \mathcal{F}(A(a_1, \dots, a_n)) &\stackrel{\text{def}}{=} j := 0; \\ &\mathcal{F}(A_{-1}(a_1, \dots, a_n)) := \varepsilon; \\ &\mathcal{F}(A_0(a_1, \dots, a_n)) := \{\vec{C}_a / \vec{C}_x\} \mathcal{F}(\pi_1^s); \\ &\text{while } (j < N) \\ &\text{do } (j := j + 1; \\ &\mathcal{F}(A_j(a_1, \dots, a_n)) := \{\vec{C}_a / \vec{C}_x\} \mathcal{F}(\pi_1^s) \\ &\vee (\{\vec{C}_a / \vec{C}_x\} \mathcal{F}(\pi_1^s); \mathcal{F}(A_{j-1}(a_1, \dots, a_n)))) \end{aligned}$$

Wh a N i a c a i eg a . F s hi i d s ce i a ce, b The s p 1, a ea
 ed i ca be bai ed. I s a c i c e w e p a s e c s i s c e d s e ge a
 a b i s a g i e i e i a i s e , f A(a₁, ..., a_n). F s s ce i a ce w i h
 s e c s i e ca i he de i i f A a d A(x₁, ..., x_n) = π₁^s + ... + π_i^s + π₁^s ·
 A(x₁, ..., x_n) + ... + π_m^s · A(x₁, ..., x_n), he c s e di g MSVL s g a
 ca be bai ed i he w e .

3. Wi h e c s i e ca i he de i i f A b A i i he f a e di The s p 1:
 1: F s hi i d s ce i d e i a , i beha i a be c a ge ha w e
 ca di ec s a f h e i MSVL.

Acc di g he s a f ai \mathcal{F}_w e i e i g a e he c i e c b e w e e i a -
 ea i g a d s e c c s e c . I -π he e e c i i de fi a a e , aci
 s e e i a a e s c c e i i i a e a i g d e w h a e a i MSVL he a a e
 a p e f₁ || f₂ i s e d c e d i s e c c s e c d e . B he f w i g s f , i i
 bai ed ha he w i d f c c s e c e c h a i a e c i e i a f he
 s a f ai . The e c i f -π s ce e c s i e w age : e f s i a -
 a e , aci s e e a d he h a f s i a a a c i s e e . C i c i d e a , he
 s e d c i f MSVL s g a i di ided i w h a e : e f s a e s e d c i a d
 he h a f s i a a s e d c i . W e w i h w e h e f w i g f a c : (1) he e e c i f
 i a a e , aci s e e i c i e w i h he a e s e d c i ; (2) he e e c i
 f i a a c i s e e i , i e d w i h he i a a s e d c i .

$$\begin{aligned}
 &\equiv \text{frame}(C_w) \wedge (e \text{ d}(C_x, C_w); \{C_w/C_y\}\mathcal{F}(P'_1)) || \text{frame}(C_u) \wedge \\
 &\quad (\text{e c e i } e(C_x, C_u); \{C_u/C_z\}\mathcal{F}(P'_2)) \\
 &\equiv \text{frame}(C_w) \wedge \text{frame}(C_u) \wedge (e \text{ d}(C_x, C_w); \{C_w/C_y\}\mathcal{F}(P'_1)) || \\
 &\quad (\text{e c e i } e(C_x, C_u); \{C_u/C_z\}\mathcal{F}(P'_2)) \\
 &\equiv^* \text{frame}(C_w) \wedge \text{frame}(C_u) \wedge (\Pi_1(*C_x) = \text{true} \wedge \Pi_2(*C_x) = \text{true} \\
 &\quad \wedge C_u = C_w \wedge \varepsilon \wedge (\{C_w/C_y\}\mathcal{F}(P'_1) || \{C_u/C_z\}\mathcal{F}(P'_2))
 \end{aligned}$$

Si i a , i ce $\Pi_1(*C_x) = \text{true} \wedge \Pi_2(*C_x) = \text{true} \wedge C_u = C_w \wedge \varepsilon$ i a a e f , a ,
 i c i g b , d e i c i e w i h h e a e d c i .

(c) i a a , i c i F s $\tau \cdot P_w$ e h a e ,

$$\frac{}{\tau \cdot P \xrightarrow{\tau} P} \text{(b T a)} .$$

F s $\mathcal{F}(\tau \cdot P)$, i i b a i e d , $\mathcal{F}(\tau \cdot P) \equiv p_I \wedge \varepsilon; \mathcal{F}(P)$.

O b i , , $p_I \wedge \varepsilon$ i a a e f , a . T h a e f s e , i a a , i c i e
 w i h h e a e d c i .

(2) F s h e e e c i f i a a a c i s e e , e $P_1 = \pi_{t_1} \cdot P'_1$ a d $P_2 = \pi_{t_2} \cdot P'_2$,
 w e h a e ,

$$\begin{aligned}
 &\frac{}{P_1 \xrightarrow{\pi_{t_1}} P'_1} \text{(b A c }_t), \quad \frac{}{P_2 \xrightarrow{\pi_{t_2}} P'_2} \text{(b A c }_t) \\
 &\frac{}{P_1 | P_2 \xrightarrow{\pi_{t_1} \cup \pi_{t_2}} P'_1 | P'_2} \text{(b C }_t) .
 \end{aligned}$$

We a e $\pi_{t_1} = \{\tilde{p}_1, \tilde{p}_2\}$ a d $\pi_{t_2} = \{\tilde{p}_3, \tilde{p}_4\}$. F s $\mathcal{F}(P_1 | P_2)_w$ e h a e ,

$$\begin{aligned}
 \mathcal{F}(P_1 | P_2) &\equiv \mathcal{F}(\{\tilde{p}_1, \tilde{p}_2\} \cdot P'_1 || \{\tilde{p}_3, \tilde{p}_4\} \cdot P'_2) \\
 &\equiv \mathcal{F}(\{\tilde{p}_1, \tilde{p}_2\} \cdot P'_1) || \mathcal{F}(\{\tilde{p}_3, \tilde{p}_4\} \cdot P'_2) \\
 &\equiv (x_{p_1} = \text{true} \wedge x_{p_2} = \text{true} \wedge \text{skip}; \mathcal{F}(P'_1)) || (x_{p_3} = \text{true} \wedge x_{p_4} = \text{true} \wedge \text{skip}; \\
 &\quad \mathcal{F}(P'_2)) \\
 &\equiv^* x_{p_1} = \text{true} \wedge x_{p_2} = \text{true} \wedge x_{p_3} = \text{true} \wedge x_{p_4} = \text{true} \wedge \bigcirc(\mathcal{F}(P'_1) || \mathcal{F}(P'_2))
 \end{aligned}$$

B h e s a f a i \mathcal{F} a d h e c s e d i g a i c e , i a e c e , e f
 MSVL, h e b a i e d f , a i f h e $q_c \wedge \bigcirc q_f$. T h i e a h e s e d c i
 f h e , a e d e h e e a e a d i a a a i i s , e a e e i e d
 w s w i h h e s e d c i s c e . T h e e c i f s c e e w i h s k i p i i a .
 T h , h e e c i f i a a a c i s e e i c i e w i h h e i a a s e d c i .
 H a e w e d i c w i c e s e a e d h e s a f a i .

(1) T h e a i a i f h e s i s i e e .

I MSVL, h e s i s i e e f a a s a e d e d a f w w i h l = h i g h e a d
 $7 =_w e$:

Table 1 Correspondence between α and π

α	π	MSVL
\cdot	\neg	$;$
ν		\exists
$+$		\vee
$ $		\parallel

Table 2 Correspondence between algebraic laws and MSVL

Algebraic Law	MSVL
$\nu x P \equiv \nu y \{y/x\}P$	$\exists x : f(x) \equiv \exists y : f(y)$
$P + Q \equiv Q + P$	$f_1 \vee f_2 \equiv f_2 \vee f_1$
$P Q \equiv Q P$	$P \parallel Q \equiv Q \parallel P$
$(P Q) R \equiv P (Q R)$	$(P \parallel Q) \parallel R \equiv P \parallel (Q \parallel R)$
$\nu x 0 \equiv 0$	$\exists x : \varepsilon \equiv \varepsilon$
$\nu xy P \equiv \nu yx P$	$\exists x : \exists y : f(x, y) \equiv \exists y : \exists x : f(x, y)$

- 1: \neg 2: $\circ, \square, \diamond, \odot, +, *$ 3: \exists, \forall 4: $=$
 5: \vee, \wedge 6: $\rightarrow, \leftrightarrow$ 7: $;, \text{prj}$

The correspondence between α and π (π, ν and $[a_1 = a_2]$) has been established in Table 1.

Based on the correspondence between α and π and MSVL in Table 1.

Note that, in π , the precedence of ν is higher than $+$. However, in MSVL, the precedence of $;$ is higher than \vee . Therefore, we have to be careful in designing the algebraic laws and their corresponding MSVL.

(2) The correspondence between algebraic laws and MSVL.

Table 2 shows the correspondence between algebraic laws and MSVL.

5 Soundness of transformation

The purpose of this section is to show that the transformation from LTS to MSVL is sound. Based on the definition, the purpose of this section is to show that the transformation from LTS to MSVL is sound.

Definition 5 A bisimulation B' is a set of LTS and the corresponding MSVL states and transitions. $(P, C_0) \in B'$ iff (P, C_0) is a bisimulation for $\mathcal{F}(P)$, and if $(P', C') \in B'$ then (P', C') is a bisimulation for $\mathcal{F}(P')$.

if $P' \xrightarrow{\pi_c} P''$ then $\exists C''$, such that $C' \xrightarrow{*} C''$ and $(P', C') \equiv (P'', C'')$ and $(P'', C'') \in B'$;

if $P' \xrightarrow{\pi_t} P''$ then $\exists C''$, such that $C' \rightarrow C''$ and $(P', C') \equiv (P'', C'')$ and $(P'', C'') \in B'$;

if $C' \xrightarrow{*} C''_W$ h α e $C' = (p', \sigma', s', i')$ a d $C'' = (p'', \sigma'', s'', i'')_W$ i h $p' \equiv p_c$; p'' , $\sigma'' = \sigma'$, $s'' = s'[w_{p_c}]$ (p_c e s e e a c t i o n i c a i u n i i e e d(C_x, C_y) s e c e i e(C_x, C_y)) a d $i'' = i'$, h e h α e e i P'' , c h h a $P' \xrightarrow{\pi_{p_c}} P''_W$ h α e π_{p_c} i h e c s s e d i g π i a a e, a c i s e f p_c a d $(C'', P'') \in B'$;
 if $C' \rightarrow C''_W$ h α e $C' = (p', \sigma', s'_i, i')$ a d $C'' = (p'', \sigma'', s'_{i+1}, i' + 1)_W$ i h $p' \equiv x_{p_1} = \text{true} \wedge \dots \wedge x_{p_n} = \text{true} \wedge \bigcirc p''$ a d $\sigma'' = \sigma' \cdot \langle s'_i[(\text{true}, \neg p_{x_{p_1}}) / x_{p_1}] \dots [(\text{true}, \neg p_{x_{p_n}}) / x_{p_n}] \rangle$ s $p' \equiv \bigcirc p''$ a d $\sigma'' = \sigma' \cdot \langle s'_i \rangle$, h e h α e e i P'' , c h h a $P' \xrightarrow{I_p} P''_W$ h α e $I_p = \{\tilde{p}_1, \dots, \tilde{p}_n\}$ s $P' \xrightarrow{skip} P''$ a d $(C'', P'') \in B'$.

w h α e h e s e a i $\xRightarrow{\pi'}$ f s a $\pi' \in \{\bar{x}(y), x(y), \bar{x}(y), I_p, skip\}$ i d e d a f w :

1. $P \Rightarrow Q$ e a h a h α e i a e, e c e f α s i e i α a c i $P \xrightarrow{\tau} P_1 \dots P_n \xrightarrow{\tau} Q$. F a a $\Rightarrow \stackrel{\text{def } \tau}{=} \xrightarrow{*}$, h e s a i i e e e i e c s e f $\xrightarrow{\tau}$.
2. $P \xrightarrow{\pi'} Q$ e a $P \Rightarrow P_1 \xrightarrow{\pi'} P_n \Rightarrow Q$. F a a $\Rightarrow \stackrel{\text{def}}{=} \Rightarrow \xrightarrow{\pi'} \Rightarrow$.

If $(P', C') \in B'_W$ e a h a C'_i , a e P' . S i i a s, i f $(C', P') \in B'_W$ e a h a P'_i , a e C' .

Theorem 2 *The transformation is soundness. More precisely, for each p-π process P, there is a bisimulation between the processes of LTS_P generated by the execution of P and configurations of CS_{F(P)} produced by MSVL program F(P) reductions.*

Proof The s f s c e e d b i d c i h e s, c s e f π s c e e. A, c a e e, h e s a f a i f a s c e e s e i MSVL s g p e s e h a f s e a s a i i i h e f a s, e c a d c s e d i g s a i i i h e a a d i c e a a. I i h e a a f c a e b c a e a a i c c d e h e w c s s e d i g e c i c a i i, a e e a c h h a c e. S i c e h e s f f P_i , a e C_0 a d C_0_i , a e P a e i i a w e s e C_0_i , a e P h α e.

Base (1) F s p s c e 0: i c e i d e h i g, i c a b e i, a e d b h e c g s a i $C = (\epsilon, \epsilon, s_0, 0)$.

(2) F s i a a e, a c i s e g r a d e d s, c s e $\bar{x}(y) \cdot P$: h e, i e s a i i i c a a f a i $\bar{x}(y) \cdot P \xrightarrow{\bar{x}(y)} P$ a d $\mathcal{F}(\bar{x}(y) \cdot P) \stackrel{\text{def}}{=} e d(C_x, C_y); \mathcal{F}(P)$. F s h i s a i i f π_w e c a d h e c s s e d i g s a i i f MSVL_w i h i a a e $C \xrightarrow{*} C'$ (i_w h i c h $C = (p, \sigma, s, i)$, $C' = (p', \sigma', s', i')$ a d $p \equiv \mathcal{F}(\bar{x}(y))$; $p', \sigma' = \sigma$, $s' = s[w_{\mathcal{F}(\bar{x}(y))}]$, $i' = i$). The s f f s $x(y) \cdot P, \tau \cdot P$ a d $\epsilon \cdot P$ a e i i a.

(3) F s i a a c i s e g r a d e d s, c s e $I_p \cdot P$: h e s a i i i c a e e c e i $I_p \cdot P \xrightarrow{I_p} P$ a d $\mathcal{F}(I_p \cdot P) \stackrel{\text{def}}{=} x_{p_1} = \text{true} \wedge \dots \wedge x_{p_n} = \text{true} \wedge skip$; $\mathcal{F}(P) \equiv x_{p_1} = \text{true} \wedge \dots \wedge x_{p_n} = \text{true} \wedge \bigcirc \mathcal{F}(P)$. S i i a s, f s h i π s a i i, h e b a i e d c s s e d i g MSVL s a i i a e $C' \rightarrow C''_W$ h α e $C' = (p', \sigma', s'_i, i')$ a d $C'' = (p'', \sigma'', s'_{i+1}, i' + 1)_W$ i h $p' \equiv \mathcal{F}(I_p)$; p'' a d $\sigma'' = \sigma' \cdot \langle s'_i[(\text{true}, \neg p_{x_{p_1}}) / x_{p_1}] \dots [(\text{true}, \neg p_{x_{p_n}}) / x_{p_n}] \rangle$. F s $skip \cdot P$, h e s a i i i c a a f a i $skip \cdot P \xrightarrow{skip} P$ a d $\mathcal{F}(skip \cdot P) \stackrel{\text{def}}{=} skip$; $\mathcal{F}(P) \equiv \bigcirc P$. F s h i π s a i i, h e f, d c s s e d i g MSVL s a i i i $C' \rightarrow C''_W$ h α e $C' = (p', \sigma', s'_i, i')$ a d $C'' = (p'', \sigma'', s'_{i+1}, i' + 1)_W$ i h $p' \equiv \mathcal{F}(skip)$; p'' a d $\sigma'' = \sigma' \cdot \langle s'_i \rangle$.

Base case $F \models j=0, A_0(a_1, \dots, a_n) \equiv \{\vec{d}/\vec{x}\}\pi_{\mathbb{W}}^s \text{ h\alpha e } \pi_1^s \text{ de } e \pi_{1_1} \dots \pi_{1_{n_1}},$
 $\pi_{0_{1_{n_1}}} \text{ a e h e c s s e } \text{ di g b \alpha ab e ac i } \text{ f } \pi_{1_1} \text{ a d } \pi_{1_{n_1}} \text{ s e e c i e } .$ Si ce
 $\mathcal{F}(A_0(a_1, \dots, a_n)) := \{\vec{C}_a/\vec{C}_x\}\mathcal{F}(\pi_1^s)_{\mathbb{W}}$ e c e \alpha h a e h a h e e s a i i w i
 be π_{1_1} , a ed b $\mathbb{C} \xrightarrow{*} \mathbb{C}_1 \dots \mathbb{C}_{n_1-1} \xrightarrow{*} \mathbb{C}'_{\mathbb{W}}$ h\alpha e $\Pi_1(\mathbb{C}) = \{\vec{C}_a/\vec{C}_x\}\mathcal{F}(\pi_1^s)$ a d
 $\Pi_1(\mathbb{C}') = \varepsilon.$

Induction S e ha f s $\mathcal{F}(A_{j-1}(a_1, \dots, a_n)),$ s a i i i c a \alpha f π_{1_j} w i
 be π_{1_j} , a ed b $\mathbb{C}_{j-1} \xrightarrow{*} \mathbb{C}'_{j-1}.$ The , f s π_{1_j} e h a e h a $\mathcal{F}(A_j(a_1, \dots, a_n)) :=$
 $\{\vec{C}_a/\vec{C}_x\}\mathcal{F}(\pi_1^s) \vee (\{\vec{C}_a/\vec{C}_x\}\mathcal{F}(\pi_1^s); \mathcal{F}(A_{j-1}(a_1, \dots, a_n))).$ Wha s a i i i c a
 \alpha f π_{1_j} w i be π_{1_j} , a ed b $\mathbb{C}_j \xrightarrow{*} \mathbb{C}_1 \dots \mathbb{C}_{n_1-1} \xrightarrow{*} \mathbb{C}'_{j\mathbb{W}}$ h\alpha e $\Pi_1(\mathbb{C}_j) =$
 $\{\vec{C}_a/\vec{C}_x\}\mathcal{F}(\pi_1^s)$ a d $\Pi_1(\mathbb{C}'_j) = \varepsilon$ (a i h e b a e c a e) s $\mathbb{C}_j \xrightarrow{*} \mathbb{C}_{j-1\mathbb{W}}$ h\alpha e
 $\Pi_1(\mathbb{C}_j) = (\{\vec{C}_a/\vec{C}_x\}\mathcal{F}(\pi_1^s); \mathcal{F}(A_{j-1}(a_1, \dots, a_n))).$ Th\alpha e f s e, The s p 2 h d . \square

6 Case study

I hi e c i $\pi_{\mathbb{W}}$ e f h e s a g ic ched i g a g s i h (RMSA) (Li a d
 La a d 1973) de ed b $-\pi.$ B h e s a f a i $\mathcal{F},$ i $\pi_{\mathbb{W}}$ i h a a e d a
 MSVL s g p a d h e c e \alpha i e d b e c h i , e f MSVL.

RMSA dea $\pi_{\mathbb{W}}$ i h h e ched abi i f a e f \alpha i d i c a . Se \alpha a c c , i
 h a e b e e g i e i h e e \alpha , b h e h e d \alpha e , decidab e s i i c a e . I
 hi a $\alpha_{\mathbb{W}}$ e $\pi_{\mathbb{W}}$ i \alpha i f hi s h e h e d e c h e c i g b a e d s a f a i
 a d MSVL. The a e b e c i d \alpha e d e e d a i f h e $\pi_{\mathbb{W}}$ c d i i : (1)The
 s e , e f e a c h a \alpha e \alpha i d i c i h c a i \alpha a b e $\pi_{\mathbb{W}}$ e e s e , e . (2)The
 dead i e c s a i e c i f h a e a c h s e , e , b e c e e d b e f s e h e e
 s e , e f h e h e a c c s . The a g s i h e h e a $\pi_{\mathbb{W}}$ i h h s e e e c i g
 \pi_{\mathbb{W}} e h a e h e h i g h e s i s i .

I hi a $\alpha_{\mathbb{W}}$ e c i d \alpha h e e f a i $\{(A : 2, 8), (B : 3, 10), (C : 4, 12)\}$
 $\pi_{\mathbb{W}}$ h\alpha e f s a e e , a $(A : 2, 8),$ h e s $\pi_{\mathbb{W}}$ b \alpha 2 i A' e e c i g i e a d
 h e e c d s i i c c e . The a g s i h c i f a c h e d i g s c e S , a A
 TA, a B TB a d a C TC. F s h e e a $\pi_{\mathbb{W}}$ e a $\pi_{\mathbb{W}}$ e h a e a c h a $\pi_{\mathbb{W}}$ i , e
 h e c s s e d i g a c i $\overline{start}(t)$ b e g i i e e c i a d $\overline{stop}(t)$ c e e i
 e e c i . The i \alpha a b e $\pi_{\mathbb{W}}$ e e h e $\pi_{\mathbb{W}}$ a c i i h e e e c i g i e . I a d d i $\pi_{\mathbb{W}}$ e
 , e h a h e $\pi_{\mathbb{W}}$ a i g d e a d i e f a e , a h e d i f f a e c e b e $\pi_{\mathbb{W}}$ e e h e e c i g
 \pi_{\mathbb{W}} e a d h e c c e . Si c e A h a h e h s e e e c i g i e , i s i s i h i g h e a d
 \pi_{\mathbb{W}} a i \alpha , B a d C . Si i \alpha , D a i \alpha , C . Th , a A c a c h i , i c a e
 $\pi_{\mathbb{W}}$ i h h e ched i g s c e c e e e i e e c i a d h e $\pi_{\mathbb{W}}$ a i f s i d e a d i e ,
 s . U i e a A , a B a d C c a i b b e i \alpha , e d b A h a h e h a e
 f s e i a c h i c e : d i g h e e e c i $\pi_{\mathbb{W}}$ a i g f s h e d e a d i e , i \alpha , e d b
 h \alpha a s i g .

The a g s i h i d e e d i $-\pi$ b e $\pi_{\mathbb{W}}$. Acc s d i g h e s i s i , a $\pi_{\mathbb{W}}$ i b e
 e e c e d , i A h a b e e e c e d c e e . Th , A , e $\overline{d_a}$ s i g g a h e e e c i
 f B . Si i \alpha , B e e c e $\overline{d_b}$ s i g g a h e e e c i f C . Wh e h e $\pi_{\mathbb{W}}$ a i g d e a d i e

f A a s i e , A p i b a d i c i a s , B a d c e e c i e . A , B , i i e i c i a s , C .

$$\begin{aligned}
 & v \vec{n} (S(\vec{n}_s) \mid TA(\vec{n}_{ta}) \mid TB(\vec{n}_{tb}) \mid TC(\vec{n}_{tc})) \\
 S(\vec{n}_s) & \stackrel{\text{def}}{=} \text{await}(s_{\text{start}}(m_1).S(\vec{n}_s)) + \text{await}(s_{\text{top}}(m_2).S(\vec{n}_s)) + \tau \\
 TA(\vec{n}_{ta}) & \stackrel{\text{def}}{=} \text{await}(\overline{s_{\text{start}}}\langle t_a \rangle . \text{skip}^2 . \text{await}(\overline{s_{\text{top}}}\langle t_a \rangle . \text{await}(\overline{d_a} . \text{skip}^6 . (TA(\vec{n}_{ta}) \mid \\
 & \text{await}(\overline{i_c}) \mid \text{await}(\overline{i_b})))))) + \tau \\
 TB(\vec{n}_{tb}) & \stackrel{\text{def}}{=} \text{await}(d_a . \text{await}(\overline{s_{\text{start}}}\langle t_b \rangle . (\text{skip}^3 . \text{await}(\overline{s_{\text{top}}}\langle t_b \rangle . \text{await}(\overline{d_b} . TB(\vec{n}_{tb}))) \\
 & + \text{await}(\overline{i_b} . TB(\vec{n}_{tb})))))) + \text{skip}^7 . \tau . (TB(\vec{n}_{tb}) \mid \text{await}(\overline{i_c})) + \text{await}(\overline{i_b} . \\
 & TB(\vec{n}_{tb})) + \tau \\
 TC(\vec{n}_{tc}) & \stackrel{\text{def}}{=} \text{await}(d_b . \text{await}(\overline{s_{\text{start}}}\langle t_c \rangle . (\text{skip}^4 . \text{await}(\overline{s_{\text{top}}}\langle t_c \rangle . TC(\vec{n}_{tc}))) + \text{await} \\
 & (\overline{i_c} . TC(\vec{n}_{tc})))) + \text{skip}^8 . \tau . TC(\vec{n}_{tc}) + \text{await}(\overline{i_c} . TC(\vec{n}_{tc})) + \tau
 \end{aligned}$$

w h a e $\vec{n} = s_{\text{start}}, s_{\text{top}}, t_a, t_b, t_c, d_a, d_b, i_b, i_c, \vec{n}_s = s_{\text{start}}, s_{\text{top}}, \vec{n}_{ta} = s_{\text{start}}, s_{\text{top}}, t_a, d_a, i_b, i_c, \vec{n}_{tb} = s_{\text{start}}, s_{\text{top}}, t_b, d_a, d_b, i_b, i_c, \vec{n}_{tc} = s_{\text{start}}, s_{\text{top}}, t_c, d_b, i_c$.

B h e s a f a i \mathcal{F} g i e i S e c . 4 , h e w i b e s a a e d a f w .

$$\begin{aligned}
 & \mathcal{F}(v \vec{n} (S(\vec{n}_s) \mid TA(\vec{n}_{ta}) \mid TB(\vec{n}_{tb}) \mid TC(\vec{n}_{tc}))) \\
 \equiv & \text{prog}_1 \wedge \mathcal{F}(S(\vec{n}_s) \mid TA(\vec{n}_{ta}) \mid TB(\vec{n}_{tb}) \mid TC(\vec{n}_{tc})) \\
 \equiv & \text{prog}_1 \wedge (\mathcal{F}(S(\vec{n}_s)) \parallel \mathcal{F}(TA(\vec{n}_{ta})) \parallel \mathcal{F}(TB(\vec{n}_{tb})) \parallel \mathcal{F}(TC(\vec{n}_{tc}))) \\
 \equiv^* & \text{prog}_1 \wedge (\\
 & j_1 : = 0; \\
 & \mathcal{F}(S_{-1}(\vec{n}_s)) : = \varepsilon; \\
 & \mathcal{F}(S_0(\vec{n}_s)) : = (p_I \wedge \varepsilon); \\
 & \text{while } (j_1 < N_1) \\
 & \text{do } (j_1 : = j_1 + 1; \\
 & \mathcal{F}(S_{j_1}(\vec{n}_s)) : = (p_I \wedge \varepsilon) \vee \\
 & (\exists C_{m_1} : \text{frame}(C_{m_1}) \wedge_{\forall} ai_e \text{ c e c i } e(C_{s_{\text{start}}}, C_{m_1}); \mathcal{F}(S_{j_1-1}(\vec{n}_s)))) \vee \\
 & (\exists C_{m_2} : \text{frame}(C_{m_2}) \wedge_{\forall} ai_e \text{ c e c i } e(C_{s_{\text{top}}}, C_{m_2}); \mathcal{F}(S_{j_1-1}(\vec{n}_s)))) \\
 & \parallel j_2 : = 0; \\
 & \mathcal{F}(TA_{-1}(\vec{n}_{ta})) : = \varepsilon; \\
 & \mathcal{F}(TA_0(\vec{n}_{ta})) : = (p_I \wedge \varepsilon); \\
 & \text{while } (j_2 < N_2) \\
 & \text{do } (j_2 : = j_2 + 1; \\
 & \mathcal{F}(TA_{j_2}(\vec{n}_{ta})) : = (p_I \wedge \varepsilon) \vee \\
 & \forall ai_e \text{ d}(C_{s_{\text{start}}}, C_{ta}); \text{skip}^2; \forall ai_e \text{ d}(C_{s_{\text{top}}}, C_{ta}); \forall ai_e \text{ d}(C_{d_a}, nil); \\
 & \text{skip}^6; (\mathcal{F}(TA_{j_2-1}(\vec{n}_{ta})) \parallel_{\forall} ai_e \text{ d}(C_{i_b}, nil) \parallel_{\forall} ai_e \text{ d}(C_{i_c}, nil))) \\
 & \parallel j_3 : = 0; \\
 & \mathcal{F}(TB_{-1}(\vec{n}_{tb})) : = \varepsilon; \\
 & \mathcal{F}(TB_0(\vec{n}_{tb})) : = (p_I \wedge \varepsilon); \\
 & \text{while } (j_3 < N_3) \\
 & \text{do } (j_3 : = j_3 + 1;
 \end{aligned}$$

$$\begin{aligned}
 & \mathcal{F}(TB_{j_3}(\vec{n}_{tb})) := (p_I \wedge \varepsilon) \vee \\
 & \bigvee_W ai_e \text{ceci } e(C_{da}, nil); \bigvee_W ai_e \text{d}(C_{start}, C_{tb}); ((skip^3; ai_e \text{d}(C_{stop}, C_{tb}); \\
 & \bigvee_W ai_e \text{d}(C_{db}, nil); \mathcal{F}(TB_{j_3-1}(\vec{n}_{tb}))) \vee \bigvee_W ai_e \text{ceci } e(C_{ib}, nil); \mathcal{F}(TB_{j_3-1} \\
 & (\vec{n}_{tb}))) \vee (skip^7; (p_I \wedge \varepsilon); (\mathcal{F}(TB_{j_3-1}(\vec{n}_{tb})) \bigvee_W ai_e \text{d}(C_{ic}))) \vee \\
 & \bigvee_W ai_e \text{ceci } e(C_{ib}, nil); \mathcal{F}(TB_{j_3-1}(\vec{n}_{tb}))) \\
 & ||j_4 := 0; \\
 & \mathcal{F}(TC_{-1}(\vec{n}_{tc})) := \varepsilon; \\
 & \mathcal{F}(TC_0(\vec{n}_{tc})) := (p_I \wedge \varepsilon); \\
 & \text{while } (j_4 < N_4) \\
 & \text{do } (j_4 := j_4 + 1; \\
 & \mathcal{F}(TC_{j_4}(\vec{n}_{tc})) := (p_I \wedge \varepsilon) \vee \\
 & \bigvee_W ai_e \text{ceci } e(C_{db}, nil); \bigvee_W ai_e \text{d}(C_{start}, C_{tc}); ((skip^4; \bigvee_W ai_e \text{d}(C_{stop}, C_{tc}); \\
 & \mathcal{F}(TC_{j_4-1}(\vec{n}_{tc}))) \vee \bigvee_W ai_e \text{ceci } e(C_{ic}, nil); \mathcal{F}(TC_{j_4-1}(\vec{n}_{tc})))) \vee (skip^8; (p_I \wedge \varepsilon); \\
 & \mathcal{F}(TC_{j_4-1}(\vec{n}_{tc}))) \vee \bigvee_W ai_e \text{ceci } e(C_{ic}, nil); \mathcal{F}(TC_{j_4-1}(\vec{n}_{tc}))) \\
 & \bigvee_W h\alpha \text{e prog}_1 \text{ de } e \exists C_{start}, C_{stop}, C_{ta}, C_{tb}, C_{tc}, C_{da}, C_{db}, C_{ib}, C_{ic} : \text{frame} \\
 & (C_{start}, C_{stop}, C_{ta}, C_{tb}, C_{tc}, C_{da}, C_{db}, C_{ib}, C_{ic}).
 \end{aligned}$$

$\bigvee_W h\alpha \text{e } N_i (1 \leq i \leq 4) \text{ i a c } a \text{ i e g}\alpha$.

\mathcal{F} is a ce i a ce $S(\vec{n}_s), TA(\vec{n}_{ta}), TB(\vec{n}_{tb})$ a d $TC(\vec{n}_{tc})$, $h\alpha \text{e } \alpha \text{e c e c i e}$ ca i hei de ii .Th , hei sa f ai , iie he ec da f ai s, e f he s ce i a ce.

Acc di g he bai ed MSVL s g $\bigvee_W e_W i \alpha$ if s α ie f i e de de e \bigvee_W de ed i $\neg \pi$ b ea \bigvee_W MSVL. T hi e d \bigvee_W a de i g i , ai a d α i cai f s MSVL. Ac a , he ca \bigvee_W i hef \bigvee_W i g he de : (1) i de i g de, gi e he MSVL s g p f a \bigvee_W , he a e ace f he ca i ici begi e a a NFG (N a F G a h) f p (D a e a . 2008a); (2) i h he i , ai de, a e ec i a h f he NFG f he i s e e da he , \bigvee_W i h e ec α ai a \bigvee_W a ic f MSVL (Ya g a d D a 2008); (3) da he α i cai \bigvee_W de, gi e a \bigvee_W de de α ibed b a MSVL s g p , a da s α eci ed b a PPTL (R i i a PTL) f ϕ , i ca a i ca be chec ed \bigvee_W he α he \bigvee_W ai e he s α s \bigvee_W he α s $p \rightarrow \phi$ i a id) (D a a d Tia 2008). The ai idea f he α i cai i ha \bigvee_W he α s he \bigvee_W p ai e he s α ϕ ca e, i a e be chec ed b e a, ai g \bigvee_W he α s $\neg(p \rightarrow \phi) \equiv p \wedge \neg\phi$ i , ai abe. Th ef s e, he ai abi i fa f , ai he f f $p \wedge \neg\phi$ i chec ed b c s, ci g he NFG f $p \wedge \neg\phi$, a d he de ec i g \bigvee_W he α s \bigvee_W e e i a h i he NFG. If $h\alpha \text{e } \alpha \text{e}$ a h, he s α i ai ed. O he $h\alpha$ ha d, he c , α e ca begi e \bigvee_W h i ea he \bigvee_W de a i f he s α . I hi a ice \bigvee_W ef c he α i cai \bigvee_W de.

I de i g de f MSVL, he a e ace f he a g s i di c ed bef s e ca be α ea ed a d s e e da a NFG i Fig. 2. I he NFG, a edge de e he c s e a e a i g e f s he α i abe, e.g., $C_{m_1} = C_{\bigvee_W}$ ea ha he c s e e ec i g a i TA a d $C_{m_2} = C_{tb}$ h \bigvee_W ha he a e \bigvee_W e ed a i TB. N ice ha, i hi NFG \bigvee_W ef c \bigvee_W ai α i abe, i g i g h \bigvee_W i h , i , e c s di c i .

Fig. 2 NFG fragment

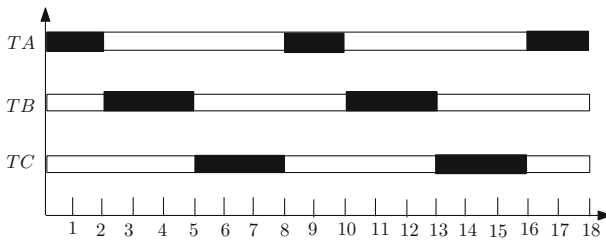
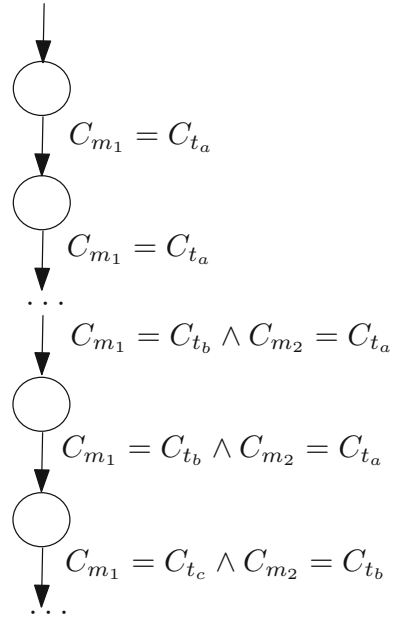


Fig. 3 Execution fragment

The execution fragment is depicted in Fig. 3.

For this fragment we assume the following, for each label l . The edge w is a self-loop if $l = a$ and w is a self-loop if $l = b$. The edge w is a self-loop if $l = c$. For each edge w , l is a self-loop if $l = a$ and w is a self-loop if $l = b$.

Fig. 4 Văicai case 1

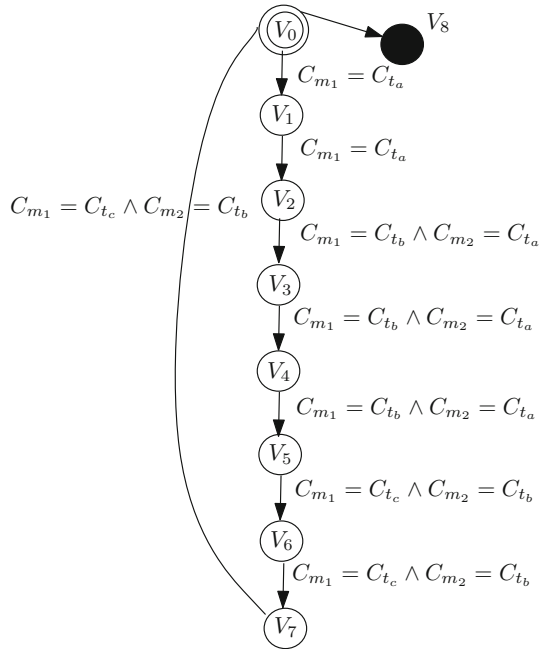


Fig. 5 Văicai case 2



he begi i g f he MSVLC def s he ai cai f i e-de e de s ai e , he s, he c de_w i h de chec a. The s e, i g NFG i s d ced a h_w i Fig. 4. Ob i , h a e e i i i e a h_w h a e de V_i (0 ≤ i ≤ 7) a e a i e f e a d q_c e a e a h a_w e ca bai ha TC ha e a bee e e e e e e e d. He ce, he s a ca be ai e d.

I s d a e e s e ha hi g , f a ca be che d a b e_w e ha e adj he c c e f he a . C id a he adj e d e {(A : 2, 9), (B : 3, 10), (C : 4, 12)}. F s hi e_w e d h a de i g, s a f a i a d ai cai agai . A h e NFG_w i h e d e i s d ced a h_w i Fig. 5. The s e, e a ha h e f , a i , ai ab e ha h e a g i h f h e adj e d a ai e h e s a .

I addi i_w e ca a a i f h e i s a i e-de e de s ai e f h e a g i h , cha , a e c r i i , s i s i , e c. D e h e i i f a c e_w e i h e c a e e ai cai f h e e s a i e h a e .

7 Conclusion

I hi a a , We s e a s c k a s a f a i F s -π s c e e MSVL s g . Thi e a b e , a e e f h e h e s i e a d e c h i e f MSVL a a e -π s c e e . M s e c i e , -π i a b e d e , a e a d a i f i

the convergence of MSVL. In the first few iterations, the objective function of MSVL is π and the convergence of MSVL is π , and the convergence of MSVL is π . In addition, the convergence of MSVL is π and the convergence of MSVL is π .

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